L24: MapReduce and DFS
Distributed Systems for the Jeff M. Phillips mashes

April 15, 2019

## Poster

## Don't put too much, or small font!

- Succinct title (and names)
- What is the problem and data you worked on?
- What were the key ideas in your approach?
- What techniques from the class did you use?
-What did you learn?
$\frac{\text { Big Data }}{\text { data too bis for }}$ Algo one computer (Stats: the limit is hire) Algo C3 count
(1) Date enormous 200-1000 machines
(2) Data Static (some appends)

Date Science Revolution
Classic Moke Mgpotlesis
Seience (2) Gather Date Eieck if dete sopports) usedne
(3) Chacte if dote sopports)
(0) Precompote, gath $i$ st Massice Date

MapReduce: (8) Google

- Boill inverted indea
- Run PaseRonk

$$
\begin{aligned}
& \text { Distribuled File System } \\
& \text { 〈tey, valec> pair JSON } \\
& \text { Zes: "onique" id } \\
& \text { volua: } \\
& \text { racy, someton } \\
& \text { messy } \\
& \text { dabe } \\
& \log \text { id } \\
& \text { wab addrest } \\
& \text { doc id } \\
& \text { wotd index } \\
& \text { actual lag } \\
& \text { htiml, out-goins } \\
& \begin{array}{c}
\text { list of masds } \\
\text { bk-grams }
\end{array} \\
& \begin{array}{l}
\text { lisd of does w/ } \\
\text { thet ward. }
\end{array}
\end{aligned}
$$

Data in Blocks

$$
\text { Block } \equiv 64 \mathrm{MB}
$$

Each flock is stored 2 or 3 times get and of locality

vern resilient (to machine failure)

Map Reduce


1. Map: Read blocker form

- poll oof import info
- determine "new licolibs" $\rightarrow$ identify new tan
1.5 Combine: Mess $\langle b, v\rangle$ before shuffle

2. Shuffle group bag bey
3. Reduce: $\overrightarrow{\rightarrow o b}$ same tree same now tho ck new bey, pood process (analysis)

Word Count (Hello World)
Consider as input all of English Wikipedia stored in DFS. Goal is to count how many times each word is used.
Map: doc id $\rightarrow\left\langle\left\langle\begin{array}{c}\text { it } \\ \text { M. }, ~ i d\rangle, ~ c w r o d, ~ i d ~\end{array}\right\rangle\right.$ $\langle$ wand, id $>$ good. 1
Combine: $\left\langle\left\langle\omega, v_{1}\right\rangle,\left\langle w, v_{2}\right\rangle, \ldots\right\rangle \rightarrow\left\langle w, \sum_{i} v_{i}\right\rangle$
 output $\left\langle\operatorname{qood}, \sum_{i} v_{i}\right\rangle$


Inverted Index
Consider as input all of English Wikipedia stored in DFS. Goal is to build an index, so each word has a list of pages it is in.
Map: artie $\rightarrow\langle$ word, id $\rightarrow$, Kurd, id $\rangle$

Reduce: (wiordir id, ${ }^{2}$, Lodi, $d_{2}$ 〉, ...

$$
\rightarrow\left\langle\operatorname{wosd}_{i}, \bigcup_{j} i_{j}\right\rangle
$$

Phrases
Consider as input all of English Wikipedia stored in DFS. Goal is to build an index, on 3 -grams (sequence of 3 words) that appears on exactly one page, with link to page.
Map : doc id $\quad \vdots\langle$ word, word e word 3, id $\rangle$ sword, wood 3 , worst I $_{1}$, id

Reduce: $\left\langle\left(w_{i}, w_{i}, w_{i n}\right), i d_{1}\right\rangle,\left(\left(w_{i} w_{1}, w_{11}\right), w_{2}\right)$

$$
\rightarrow\left\langle\left(w_{i} w_{i} w_{i \prime}\right), \bigcup_{\hat{j}} i d_{j}\right\rangle
$$

Label Propagation (Graph)
Consider a large graph $G=(V, E)$ (e.g., a social network), with a subset of notes $V^{\prime} \subset V$ with labels (e.g., \{pos, neg\}). Each node stores its label (if any) and edges.
Assign a vertex a label if (a) unlabled, (b) has $\geq 5$ labeled neighbors, (c) based on majority vote.


$$
\longrightarrow\left\langle\operatorname{lid}_{1}(t, \text { fixed })\right\rangle
$$



Reduce: $\langle i d, 1 a b\rangle,,\langle n d, \mid a b$,


Label Propagation (Embedding)
Consider a data set $X \subset \mathbb{R}^{d}$, with a subset of points $X^{\prime} \subset X$ with labels (e.g., \{pos, neg\}). Implicitly defines graph with $V=X$ and $E$ using $k=20$ nearest neighbors.
Assign a vertex a label if (a) unlabled, (b) has $\geq 5$ labeled neighbors, (c) based on majority vote.

$$
\text { LSH, in } 2 \text { rounds }
$$

Example PageRank


$$
\begin{aligned}
& q^{*}=P^{*} q . \\
& =P\left(P\left(P_{\ldots . .}\left(P_{q_{0}}\right)\right\rangle\right) \\
& \leq 50 \text { roouds }
\end{aligned}
$$

## Example PageRank

$$
\begin{aligned}
& q(1) q(6) \quad g(3) q(4)
\end{aligned}
$$

Stripes:
$M_{1}=\left[\begin{array}{c}0 \\ 1 / 3 \\ 1 / 3 \\ 1 / 3\end{array}\right] \quad M_{2}=\left[\begin{array}{c}1 / 2 \\ 0 \\ 0 \\ 1 / 2\end{array}\right] \quad M_{3}=\left[\begin{array}{l}0 \\ 1 \\ 0 \\ 0\end{array}\right] \quad M_{4}=\left[\begin{array}{c}0 \\ 1 / 2 \\ 1 / 2 \\ 0\end{array}\right]$
These are stored as $(1:(1 / 3,2),(1 / 3,3),(1 / 3,4))$, $(2:(1 / 2,1)(1 / 2,4)),\left(3:\left(1, \frac{2}{2}\right)\right)$, and $\left(4:(1 / 2,7,7),\left(1 / 2, \frac{8}{5}\right)\right)$.

## Example PageRank



Blocks:
$M_{1,1}=\left[\begin{array}{cc}0 & 1 / 2 \\ 1 / 3 & 0\end{array}\right] \quad M_{1,2}=\left[\begin{array}{cc}0 & 0 \\ 1 & 1 / 2\end{array}\right] \quad M_{2,1}=\left[\begin{array}{cc}1 / 3 & 0 \\ 1 / 3 & 1 / 2\end{array}\right] \quad M_{2,2}=\left[\begin{array}{cc}0 & 1 / 2 \\ 0 & 0\end{array}\right]$
These are stored as $(1:(1 / 2,2)),(2:(1 / 3,1))$, as
$(2:(1,3),(1 / 2,4))$, as $(3:(1 / 3,1)),(4:(1 / 3,1),(1 / 2,2))$, and as $(3:(1 / 2,4))$.
$\left\langle i d_{4}, \overline{g(1) \cdot \frac{1}{3}+g(1) \cdot \frac{1}{2}}\right\rangle$

